



Revisiting Double-Spending Attacks on the Bitcoin Blockchain: New Findings

Jian Zheng, Huawei Huang, Canlin Li, Zibin Zheng, Song Guo*
School of Computer Science and Engineering, Sun Yat-Sen University, Guangzhou, China
*Department of Computing, The Hong Kong Polytechnic University, Hong Kong.

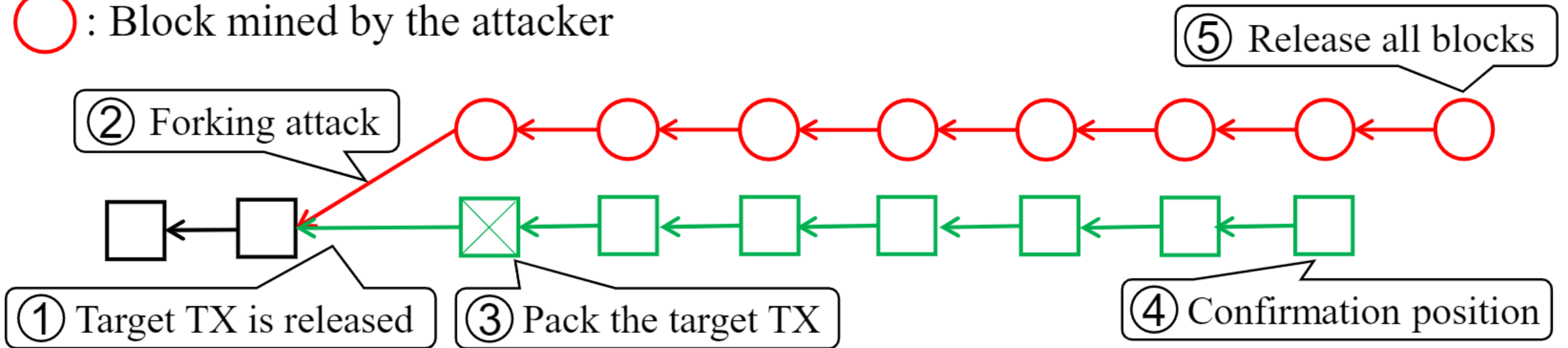
Presentation for IWQoS 2021

Double-Spending Attack



□ : Block mined by an honest miner

○ : Block mined by the attacker



Our contribution

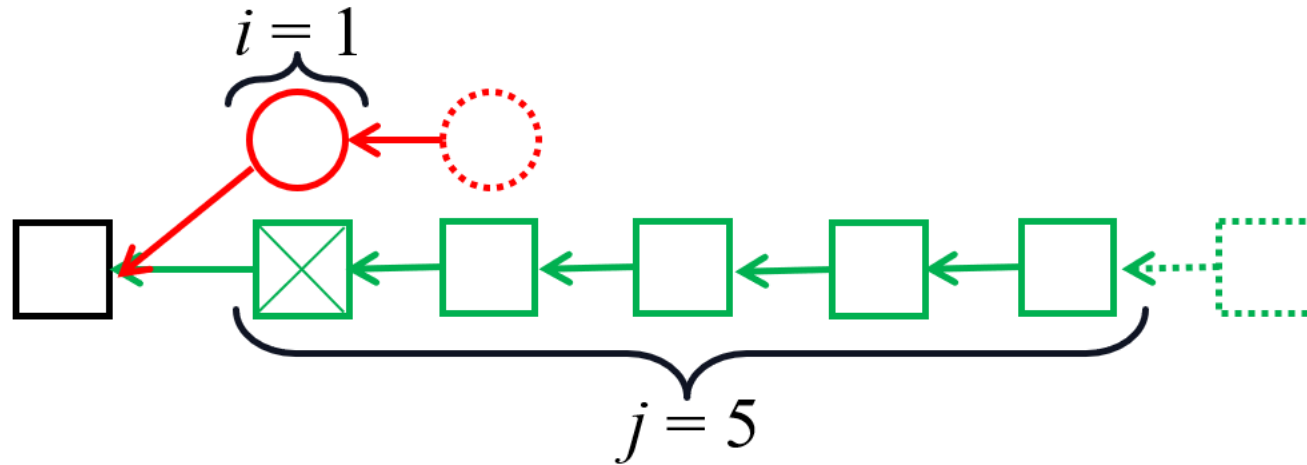


- **Adaptive DSA:** Adaptive Double-Spending Attack
- Profit-maximized attack strategy by Stochastic Dynamic Programming (SDP)
- Analytical model for simulation

An unfavourable situation to the attacker



 : The attacker is mining  : Honest miners are mining

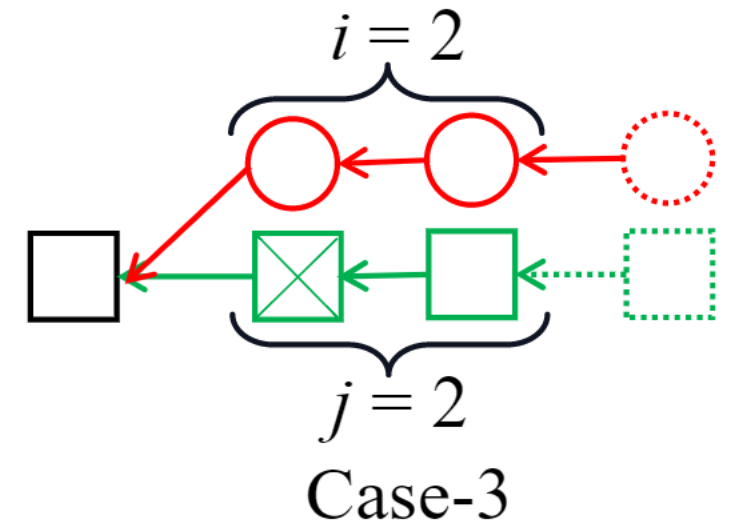
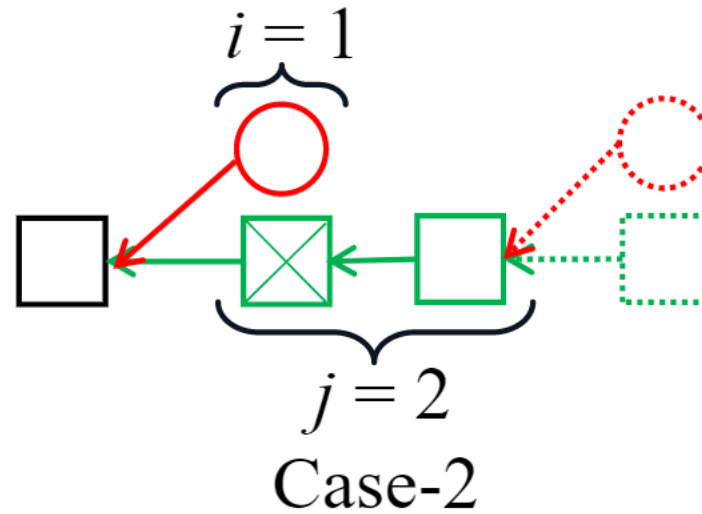
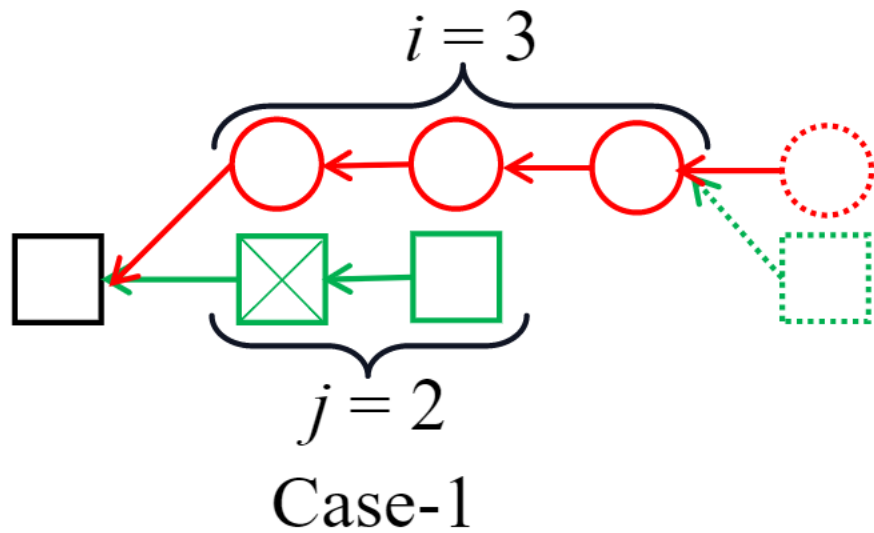


- The transaction has not been confirmed in $S_{1,5}$.
- **Should the attacker give up?**

Attack decision: Quit or Keep



- **0**: Quit attacking

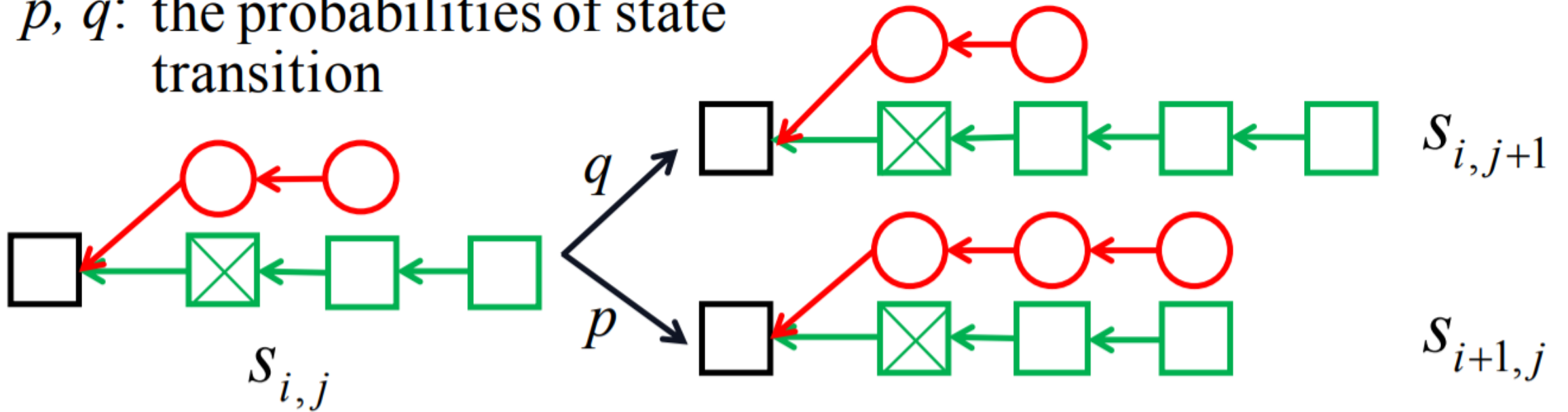


Attack decision: Quit or Keep



- 1: Keep attacking

p, q : the probabilities of state transition



Attack-Decision Matrix



• **Variables:** $d_{i,j} \in \{0, 1\}$

- **0:** Quit attacking
- **1:** Keep attacking

• **Profit:** $f(\{d_{i,j}\})$

• **Tools:**

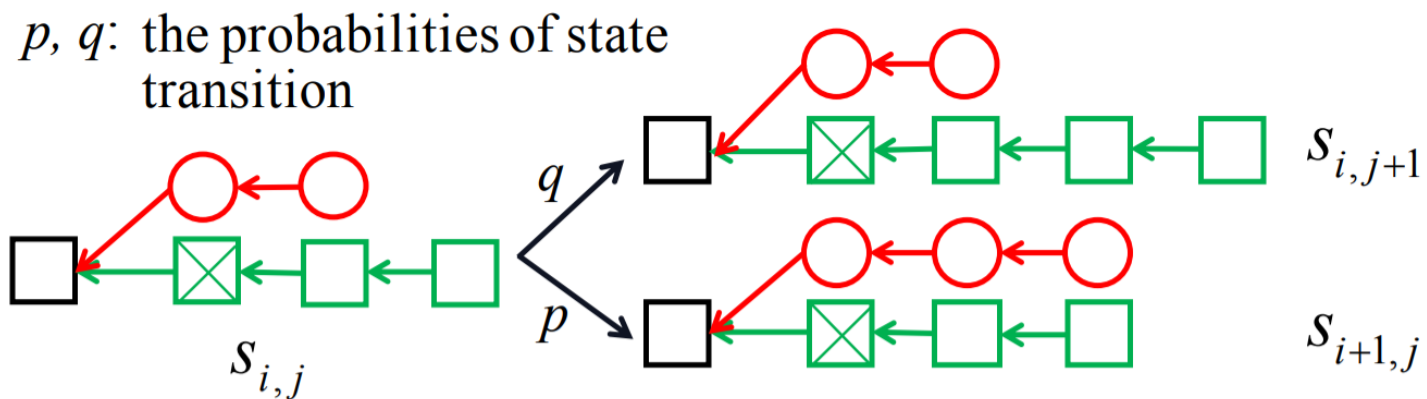
- Occurrence-Probability Matrix
- Reward Matrix

$d_{i,j}$	0	1	2	3	4	5	6
0	?	?	?	?	?	?	0
1	?	?	?	?	?	?	0
2	?	?	?	?	?	?	0
3	?	?	?	?	?	?	0
4	?	?	?	?	?	?	0
5	?	?	?	?	?	?	0
6	0	0	0	0	0	0	

Occurrence-Probability Matrix



$$P_{i,j} = \begin{cases} 1, & \text{if } i = j = 0; \\ d_{i-1,j} \cdot p \cdot P_{i-1,j}, & \text{if } j = 0 \text{ or } i = z; \\ d_{i,j-1} \cdot q \cdot P_{i,j-1}, & \text{if } i = 0 \text{ or } j = z; \\ d_{i-1,j} \cdot p \cdot P_{i-1,j} + d_{i,j-1} \cdot q \cdot P_{i,j-1}, & \text{otherwise.} \end{cases} \quad (1)$$

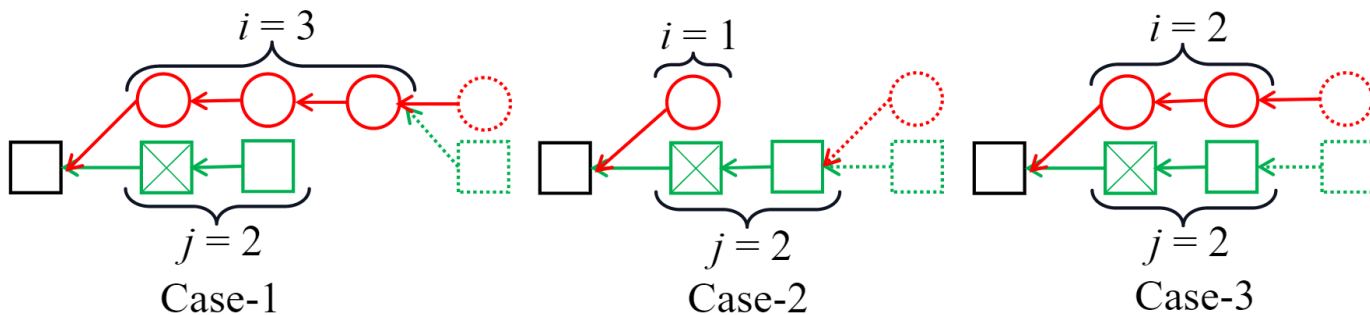


$P_{i,j}$	0	1	2	3	4	5	6
0	1	?	?	?	?	?	?
1	?	?	?	?	?	?	?
2	?	?	?	?	?	?	?
3	?	?	?	?	?	?	?
4	?	?	?	?	?	?	?
5	?	?	?	?	?	?	?
6	?	?	?	?	?	?	?

Reward Matrix



$$R_{i,j} = \begin{cases} -(i+j) \cdot cost, & \text{if } i < j; \\ i \cdot d - (i+j) \cdot cost, & \text{if } z > i > j; \\ p \cdot i \cdot d - (i+j) \cdot cost, & \text{if } z > i = j; \\ p[b + d(z+1) - (i+j+1) \cdot cost] + q \cdot R_{i,j+1}, & \text{if } i = z, j \leq z-2; \\ p[b + d(z+1) - (i+j+1) \cdot cost] + q[p(b+z \cdot d) - 2z \cdot cost], & \text{if } i = z, j = z-1. \end{cases} \quad (2)$$



$R_{i,j}$	0	1	2	3	4	5	6
0	?	?	?	?	?	?	?
1	?	?	?	?	?	?	?
2	?	?	?	?	?	?	?
3	?	?	?	?	?	?	?
4	?	?	?	?	?	?	?
5	?	?	?	?	?	?	?
6	?	?	?	?	?	?	?

Profit-Maximization Problem



- **0**: Quit attacking
- **1**: Keep attacking

$$d_{i,j} \in \{0, 1\}$$

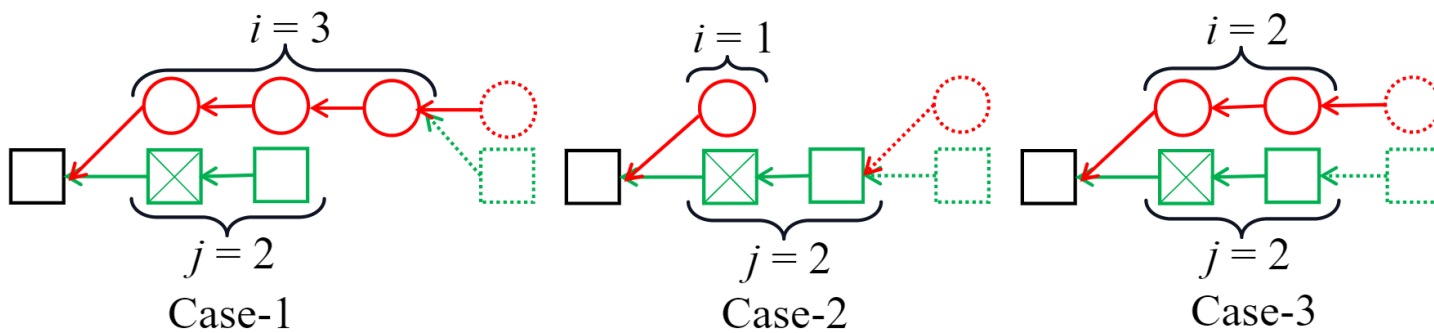
$$\max f(\{d_{i,j}\}) = \sum_{i=0}^z \sum_{j=0}^z (1 - d_{i,j}) \cdot P_{i,j} \cdot R_{i,j}$$

Stochastic Dynamic Programming (SDP)



- $d_{i,j}=0$, quit attacking

$$J_n(s_{i,j}) = J_{n+1}(s'_{i,j}) = \begin{cases} 0, & \text{if } 0 \leq i < j \leq z - 1; \\ i \cdot d, & \text{if } 0 \leq j < i \leq z - 1; \\ p \cdot i \cdot d, & \text{if } 0 \leq i = j \leq z - 1. \end{cases}$$



$J_n(s_{i,j})$	0	1	2	3	4	5	6
0	?	?	?	?	?	?	?
1	?	?	?	?	?	?	?
2	?	?	?	?	?	?	?
3	?	?	?	?	?	?	?
4	?	?	?	?	?	?	?
5	?	?	?	?	?	?	?
6	?	?	?	?	?	?	?

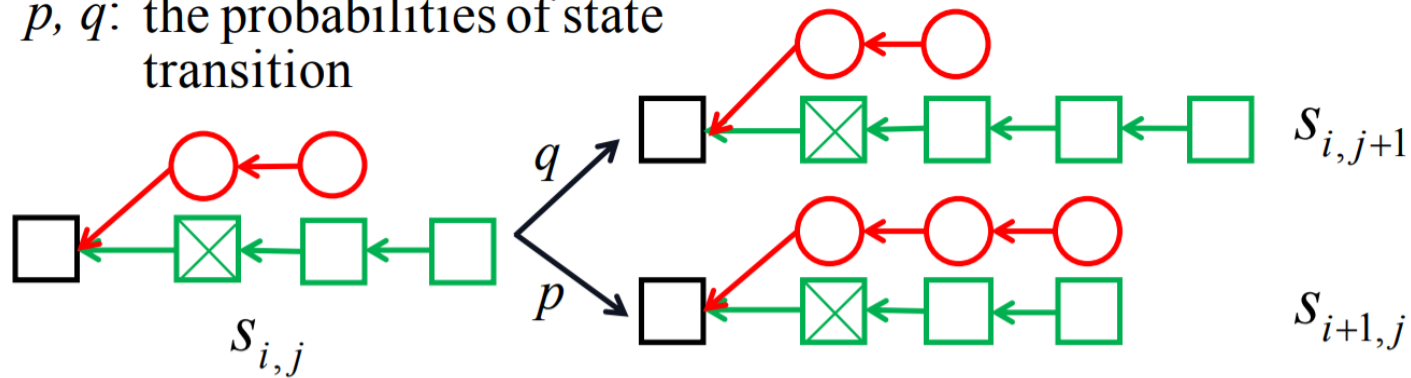
Stochastic Dynamic Programming (SDP)



- $d_{i,j}=1$, keep attacking

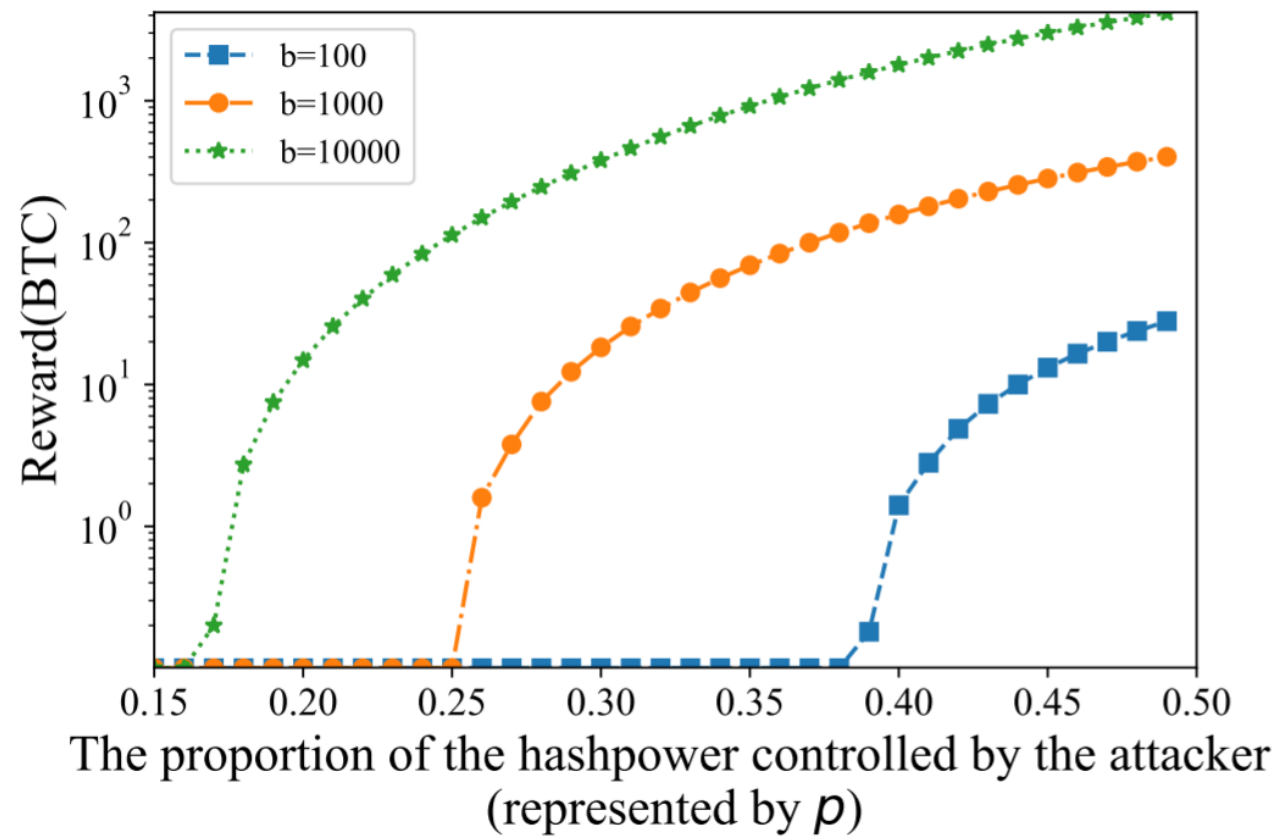
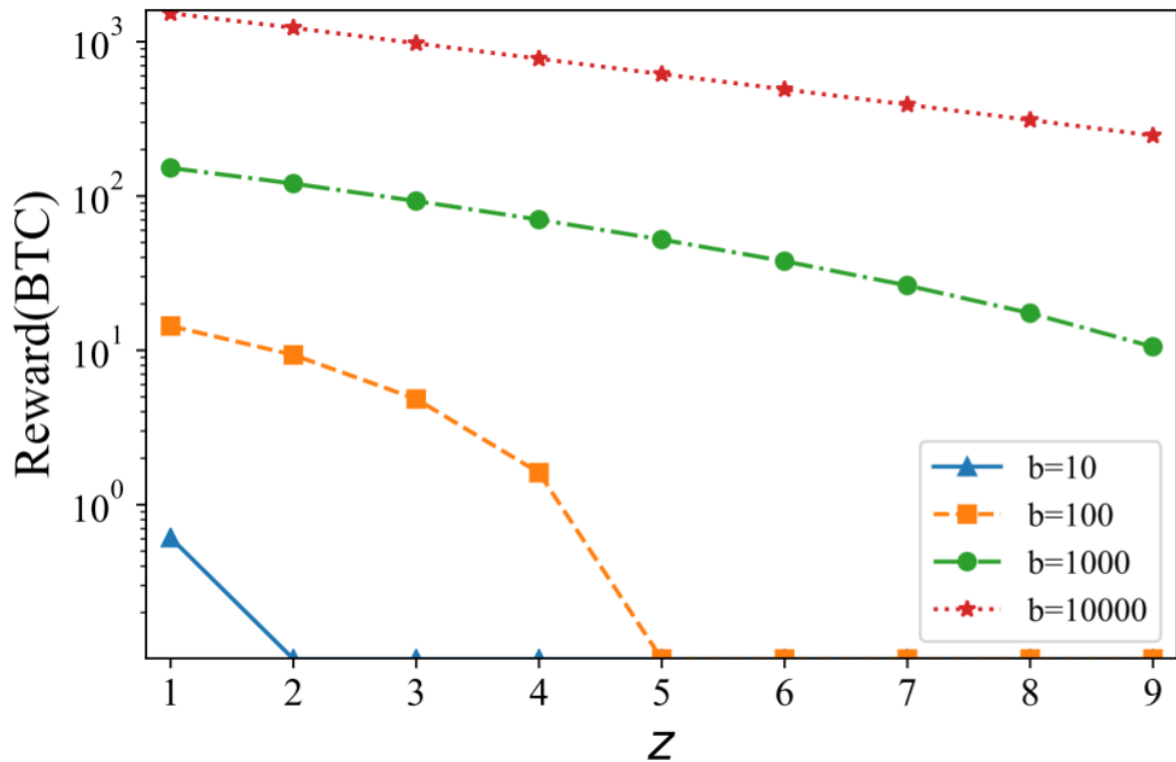
$$J_n(s_{i,j}) = -cost + p \cdot J_{n+1}(s_{i+1,j}) + q \cdot J_{n+1}(s_{i,j+1})$$

p, q : the probabilities of state transition

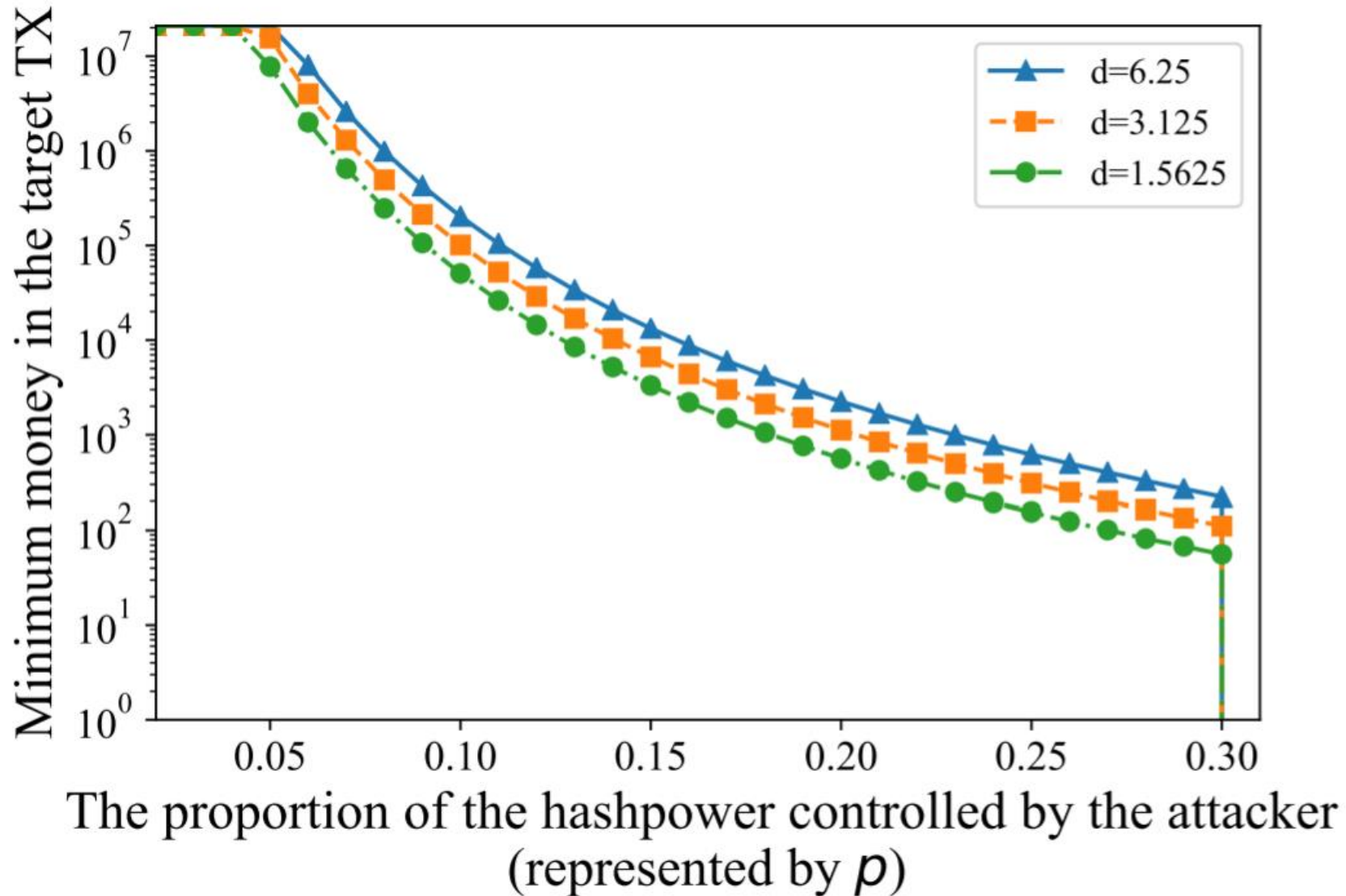


$J_n(s_{i,j})$	0	1	2	3	4	5	6
0	?	?	?	?	?	?	?
1	?	?	?	?	?	?	?
2	?	?	?	?	?	?	?
3	?	?	?	?	?	?	?
4	?	?	?	?	?	?	?
5	?	?	?	?	?	?	?
6	?	?	?	?	?	?	?

Performance Evaluation



Performance Evaluation



Conclusion



- **Adaptive DSA:** new threat to PoW-based blockchains
- 50% is not enough
- More strategies in the double-spending attack



Thanks

HuangLab@SYSU <http://xintelligence.pro/>